

U.S. PATENT APPLICATION

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Invention: SELECTABLE BANDWIDTH FACILITY FOR A NETWORK PORT

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SPECIFICATION

SELECTABLE BANDWIDTH FACILITY FOR A NETWORK PORT

Field of the Invention

5 This invention relates to data communication networks for the conveyance of data packets between users, such as computers, file servers, work stations and such like Networks of this character are commonly composed not only of the end users but also various multiport units such as hubs, switches and routers

Background of the Invention

10 Some of the units from which a network is composed are essentially passive in the sense that they are intended merely to pass on or distribute without discrimination packets that they receive Other units, particularly bridges and routers, are more complex devices and commonly include management agents by means of which control over the data packets passing through the switch may be exercised

15 In recent years many developments have occurred in the implementation of packet based networks and in particular in networks conforming to the 'Ethernet' Standards For example, the permitted maximum speeds of communication which can be achieved have increased and standards are now in existence which permit communications to occur at a variety of different speeds Devices are available which are capable of operating at a multiplicity of different rates, such as for example 10 or 100 or 1000 Mbps selectively It is known to provide on network units ports, and associated media access control devices (MACs) which are selectively controllable to operate at any one of a variety of different speeds so that a user may take advantage of a multiplicity of different operating capabilities

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30 Nevertheless, it is for a variety of reasons desirable to be able to restrict arbitrarily the bandwidth available to a user to some figure less than the rate or selected rate at which a communication link could operate For example, many network administrators do not desire their users to have access to high speed networks because allowing such access

may dramatically increase the load on existing file servers, routers and wide area network links. These reasons may be sufficient to prevent network administrators from installing multiple speed switches.

It is therefore generally desirable to provide a feature that can put an upper limit on the bandwidth permitted at an individual port of a switch. For example, a network administrator may install a switch that is capable at individual ports of operation at a multiplicity of different rates, such as 10Mbps and 100Mbps, but configure the switch such that certain ports are limited to an actual throughput of (for example) 6Mbps and other ports are limited to 20Mbps.

General State of the Art

It is known to limit the bandwidth available for packets to be transmitted over a link from a port by using a 'leaky bucket' which is decremented at a selectable rate and is incremented at a rate dependant on the traffic (i.e. the packets) dispatched from the port. Transmission of packets is inhibited if the value held in the counter exceeds a threshold. Such a system is disclosed by Harwood in United States Patent 5604867 issued 18 February 1997. It is also known to provide a system in which a peak bit rate is monitored in a first 'leaky bucket' unit and the duration of peak rate bursts is monitored in a second 'leaky bucket'. Such an arrangement is shown in Kammerl, United States patent 5339332.

It is also known to provide a 'leaky bucket' which is decremented at a selectable rate and is incremented in accordance with traffic both sent from and received by a respective port of a switch. Such a system is disclosed in GB published application GB-2336076.

It is customary in devices such as network switches, bridges and routers to store packets after they have been received and before they are transmitted from their destination port or ports. Very typically the packets are held in queues, constituted by the packet themselves or by queues of pointers for the packets, and it is well known to control the dispatch of packets from a queue either in accordance with conditions within the switch.

or by virtue of control frames received at a particular port. Accordingly, there is in most commercially available switches a mechanism which enables the transmission of switches from a particular port to be selectively inhibited.

5 It is also known to restrict the reception of packets at a particular port. In the present invention a reduction in bandwidth of packets received by the switch, is preferably achieved by inhibiting the reception of packets by signalling to a link partner, i.e. the provider of packets at the other end of a link to a particular port. This activity is well known in the art and is, for example, specified in a transmission standard such as IEEE
10 802.3 (1998 Edition) for Ethernet packets. In that standard, the sending of a particular control frame with a conventional address and particular operation code must, according to the standard, be interpreted by a receiver of such a frame (i.e. the link partner) as an instruction to cease sending frames or packets for a selectable time specified in the control frame. Such a frame (or its equivalent in other standards) will hereinafter be described as a 'pause frame'.
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Summary of the Invention

20 The present invention particularly concerns a versatile bandwidth controller by means of which at a particular port the bandwidth available for reception of packets and the forwarding of packets can be either separately or conjointly controlled. The former, called herein the duplex mode, would be more suitable when the transmission link is duplex (allowing simultaneous transmission and reception) and the latter would be more suitable when the transmission link were half-duplex (not allowing simultaneous
25 transmission and reception). For this reason the two modes will be termed 'duplex' and 'half-duplex' but it is not essential that the mode of the bandwidth controller be the same as that of the link. The link may be a duplex link but the bandwidth controller could be in a half-duplex mode if desired.

30 To this end a preferred embodiment of the invention employs two token buckets which can be configured so that the transmission and reception bandwidths may, according to the selected mode, be controlled separately or conjointly. Preferably prevention of

transmission and restriction of reception of packets may, according to the mode of the bandwidth controller, be subject to a single threshold or respective thresholds

Flow control, namely the sending of pause frames over a link and responding to pause frames, is commonly employed to relieve congestion in a switch. The sending of pause frames over a link is inherently a manner of bandwidth control. The present invention provides bandwidth control which is not dependent on congestion within a switch but can be imposed externally by a network administrator using conventional control frames

It will be understood that in the present invention, as in many other counting systems subject to a threshold, the significance of the terms 'above' and 'below' a threshold depends on the relationship (which may be arbitrary) between the threshold and the direction of counting associated with an increase in traffic. In the specific example described later, the passage of excess traffic is indicated by a net count below a threshold, but a converse convention could be employed. Stated alternatively the token buckets could be constituted by 'leaky' buckets. For this reason the term 'in-profile' is employed herein to indicate a count which will allow more traffic and 'out-of profile' is employed to indicate a count which will cause inhibition of traffic

Other features of the invention will become apparent from the following description, with reference to the accompanying drawings

Brief Description of the Drawings

Figure 1 is a general schematic illustration of a switch unit within which the present invention may be incorporated

Figure 2 is a schematic diagram illustrating one embodiment of the present invention

A typical context for the present invention is a multiport switch of the kind which is illustrated in Figure 1. The particular architecture of the switch or other unit within which the invention may be incorporated is not important. The present invention will be

described in relation to 'Ethernet' networks and more particularly those conforming to IEEE Standard 802.3, though such a restriction is not crucial to the present invention

A switch 10 as shown in Figure 1 typically has a multiplicity of ports, 11a to 11h. These ports may be in known form for coupling to transmission media. Each port is associated with a respective media access control device (MAC) 14a to 14h likewise in known form. For present purposes it may be presumed that the media access control devices conform to the requirements of IEEE Standard 802.3. They may be controllable to provide a multiplicity of transmission rates, such as 10/100/1000Mbps. Selection of a particular rate may be imposed by means of control frames or may be established by auto negotiation in a manner well known in itself (and described in the aforementioned IEEE Standard).

The remainder of the switch shown in Figure 1 is shown in schematic form only. Broadly, it comprises a bus system 13 by means of which data packets are sent for temporary storage in memory 14 and subsequently dispatched from memory to a destination port or ports. A bus system 15 between the media access control devices and processing engines 16 is provided for signalling the processing engines and also exerting control over the individual MACs. Such control is required for determining the destination of packets, and derives from both control frames received by the switch and internal operations such as the performing of lookups on address data within the packets and so on.

The basis of the system 20 shown in Figure 2 is that for a port (preferably but not necessarily each port of the switch in Figure 1), the received and transmitted streams of traffic are controlled using one or two token buckets, depending on the mode. The 'mechanism' (which may be software-based) is octet-based and can cater for bursty traffic patterns.

Preferably, there are two modes of operation. One is a full duplex mode in which two token buckets are employed. The first controls the traffic transmitted from the port by inhibiting the release from storage of packets intended for that port. The second token

bucket restricts the traffic received by the port, preferably by means of the dispatch of pause frames

A second mode of operation is a half-duplex mode. Here, a single token bucket is used to meter the combined traffic received and transmitted by the port

The system configures the token buckets to detect situations where the relevant traffic is in excess of a chosen threshold. Preferably the 'granularity' of the threshold for operation at 10 or 100Mbps is at 1Mbps intervals. For ports which are operable at 1000Mbps the granularity of the threshold could be greater, such as 10Mbps

While the traffic is below the threshold, a state termed herein 'in profile', then packets are received and/or transmitted normally without restriction by the token bucket system. If the relevant traffic exceeds the threshold, a state termed herein 'out of profile', then action is taken to stem the flow of traffic.

For the duplex mode, while the relevant token bucket indicates that the transmitted traffic is 'out of profile', then the MAC 14 is instructed to temporarily stop transmitting frames. If the MAC is instructed to stop during the transmission of a frame, then (as is usual), the frame is transmitted normally. Transmission of frames may be resumed when the token bucket indicates that the transmit traffic is once again 'in profile'

Also in the full duplex mode, while the respective token bucket indicates that the received traffic is 'out of profile', the MAC 14 is instructed to flow control incoming traffic using normal flow control methods, such as pause frames in accordance with pause frames. This flow control can be disabled when the received token bucket indicates that the traffic is once again 'in profile'

When the system is in half-duplex mode, while the single token bucket indicates that the combined traffic is 'out of profile', the MAC 14 is instructed not only to temporarily stop transmitting frames but also to flow control incoming traffic as described above

There follows a description of the elements shown in Figure 2. The MAC 14 is a normal Ethernet media access control device. The MAC 14 is coupled to a monitor 21 which, in accordance with known techniques, sends data all frames which are transmitted and received through the MAC. For each valid frame transmitted or received by the MAC,
5 the frame size (in terms of byte count) and the direction (transmitted or received) is passed to a processing engine 22. The byte count and direction may be transmitted as a data word having an appropriate operation code, one field denoting byte count and a smaller field (which may be a 1 byte field) indicating direction.

10 The processing engine 22 controls and reads configuration registers 23 (to be described) and two token buckets 24 and 25. Processing engine 22 will read the mode configuration in the configuration registers and update the appropriate token bucket. Processing engine 26 also reads the mode configuration in the configuration registers and checks whether the token buckets 24 and 25 are above or below their respective
15 counting thresholds. The threshold, which may be 'hard-wired' or set by software, defines a token count which is at the choice of the designer but would be somewhat greater than the maximum size of a packet expressed on terms of tokens. As is described later, processing engine 26 asserts signals to a receive inhibit control (RXI) 27 and a transmit inhibit control (TXI) 28.

20 Receive control 27, while the signal from processing engine 26 is asserted, inhibits the reception of more traffic by the MAC by generating a pause frame in accordance with the IEEE Standard 802.3. The duration specified in the pause frame is a matter of design choice. The transmit control, while the signal from processing engine 26 is asserted,
25 prevents the transmission of more traffic by the MAC by generating an inhibit that prevents frames being transferred from the transmit store 29 to the MAC. If the signal from engine 26 is asserted while a packet is in transit, the inhibit would be applied to the next packet.

30 In this embodiment of the invention, frames received by the MAC are stored in a receive store 30 and frames which are intended for transmission from the port are stored in a transmit storage space 29. These may be, in accordance with known practice, buffer

stores associated (temporarily or permanently) with the MAC for the particular port. Alternatively they may be allotted storage space in memory 14

Each token bucket 24 and 25 is shown explicitly in Figure 2, but may be constituted by relevant fields, as discussed below, in storage locations and be both defined and controlled by software, in accordance with the following description

Each token bucket contains a number of 'tokens', namely a specific numerical count which is related to the relevant unit of measurement of traffic flow. In this example it will be assumed that one token is equal to 1 byte. The size of the fields and thereby of the counters can be reduced if a token is equivalent to a multiplicity of bytes. This results in a loss of accuracy but for high speed measurement may be appropriate to avoid excessively large counters

In a token bucket, at regular intervals new 'tokens' are added to the bucket, specifically a 'refresh count'. A specific number of tokens are required to perform an operation, namely to send a packet or receive a packet. When the operation is performed the requisite tokens are removed from the bucket. The operation cannot be performed unless there are sufficient tokens in the bucket, i.e. the bucket count exceeds a threshold

In the present example, each token bucket is defined by four fields. These are illustrated for bucket 24 schematically as F_1 to F_4 along with a 'field' F_5 which represents the threshold. As previously mentioned, this 'field' may be a hard-wired value but could be, like fields F_1 to F_4 , set in a software controlled register

A first field, the 'in-profile' field F_1 consisting of one bit, indicates whether the token bucket is 'in profile', if the bucket count is less than a specified threshold, or whether the token bucket is out of profile

A second field F_2 , 'bucket size', having in this example three bits is a configurable value specifying the maximum number of tokens allowed in the bucket (typically values from

512 to 64k tokens) These values determine the maximum 'burst' of data which can be allowed before the token bucket becomes 'out of profile'

5 The third field F_3 is a 'refresh count', a seven bit field which indicates the number of tokens that are to be added to the bucket for each refresh interval

The refresh interval may be fixed or variable but in the present embodiment it is assumed that the refresh interval is fixed and is determined by a selected clock from the processing engine 22 to the token bucket

10 The fourth field F_4 of the token bucket is the 'bucket count' herein defined by a sixteen bit field, representing the current count of tokens in the bucket The count is reduced by incoming packets and is increased by the 'refresh count' every refresh interval, typically 8 microseconds

15 In the present example, the system includes a configuration register 23 which is controlled according to whether the mode of operation of the port is duplex or half-duplex The configuration register may be put in this mode as a result of the auto negotiation performed but it is readily possible, since the configuration registers may be set by software, to impose the configuration and the other fields on the token buckets by means of management frames sent to the port

20 Whichever scheme is adopted for the control of the configuration registers and the token buckets, the configuration register 23 in combination with processing engine 26 determines two distinct modes of operation for the token buckets, namely full duplex mode and half- duplex mode

25 In the duplex mode, the restriction of transmission of packets and the inhibition of reception of packets are separately controlled by the token buckets 24 and 25 Thus an 'in profile/out of profile' signal from token bucket 24 determines the command signal to transmit control 28 whereas the 'in profile/out of profile' signal from the token bucket 25 determines the command signal to receive control 27 More particularly, while the

'transmit' token bucket 24 indicates that the transmitted traffic is 'out of profile', control 28 prevents transfer of frames from transmit storage 29 to the port. In practice transmit control 28 instructs MAC 20 not to accept frames from transmit storage 29. As a practical matter, if the MAC is instructed to stop during the transmission of the frame, then that particular frame is transmitted normally and cessation of transmission seizes for the next frame. Transmission will be resumed when the 'transmit' token bucket indicates that the transmit traffic is once again 'in profile'.

Also in the full duplex mode, while the 'receive' token bucket indicates that the receive traffic is 'out of profile', then receive control 27 commands the MAC to produce pause frames as previously described. Flow control will be disabled when the receive token bucket indicates that the receive traffic is once again 'in profile'.

The action of the processing engine to control the signalling of controls 27 and 28 is predetermined by the configuration register 23.

For the sake of example, when the token bucket is first started, the 'bucket count' may be set to the in profile threshold (such as 0x0800 tokens). Alternatively if an initial burst of traffic is acceptable, 'bucket count' could be set to a maximum for the chosen bucket size, e.g. 0xFFFF. Every 8 microseconds, refresh tokens may be added to the bucket count. There are two options here. Option (a) requires that refresh tokens will be added to the bucket count and if the bucket count increases above the 'bucket size' the bucket count is set to the bucket size. Option (b) requires that only if the bucket count is less than the bucket size, the refresh count tokens are added to the bucket count. Either option ensures that the bucket count does not wrap around to a low number. When a packet completely arrives, the respective number of tokens is calculated. If the bucket count is greater than the in profile threshold, the calculated number of tokens is subtracted from the bucket count. If the new bucket count is equal to or less than the in profile threshold, the bucket is judged 'out of profile' and the appropriate traffic inhibiting signals will be produced. Otherwise the bucket count will not be changed.

In the full duplex mode, the bandwidth occupancy which is controlled will be the sum of the traffic to and from the port. However, for a variety of purposes that control may be too restrictive and the benefit of the present invention lies in the choice between that full duplex mode and the half-duplex mode.

In the half-duplex mode configuration register 23 sets processor 22 to cause accumulation of the byte count, irrespective of direction, in one of the token buckets, for example token bucket 24. Configuration register 23 also controls processing engine 26 so that when the threshold for token bucket 24 is exceeded, commands are sent both to transmit control 28 and receive control 27 so as to ensure that the MAC is instructed both to temporarily stop transmitting frames and also to exert flow control on the incoming traffic.

The present invention, by providing two token buckets which can control both transmitted and received bandwidth occupancy separately or as combined provides a more versatile system for scaleable control of bandwidth and in the prior art.

Degrees of sophistication are feasible. For example, it is not necessary that the thresholds in buckets 24 and 25 be the same. The configuration register could provide a choice between the token buckets to be used for the full duplex mode.